

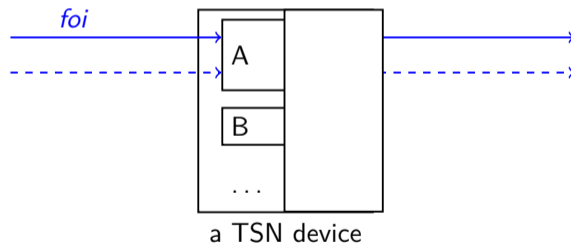
On Time Synchronization Issues in Time-Sensitive Networks with Regulators and Nonideal Clocks

published in POMACS, June 2020, doi:10.1145/3392145

Ludovic Thomas and Jean-Yves Le Boudec
ISAE-SUPAERO (France) and EPFL (Switzerland)

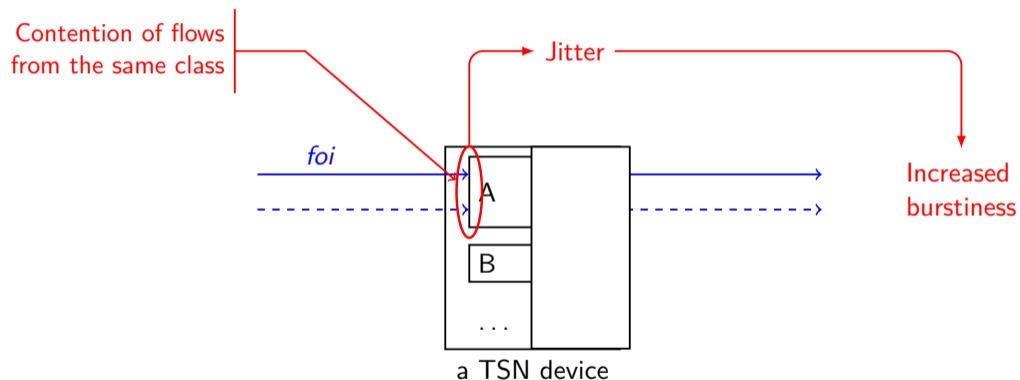
5th Workshop on Network Calculus
9 October 2020

TSN: Class-based schedulers. FIFO-per-class.



foi = flow of interest

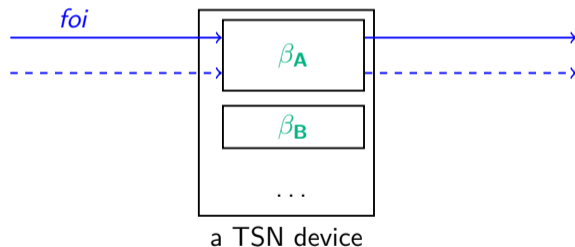
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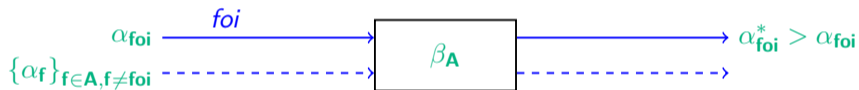
With Network Calculus (e.g. TFA [Schmitt and Zdarsky, 2006]):



TFA = Total Flow Analysis.

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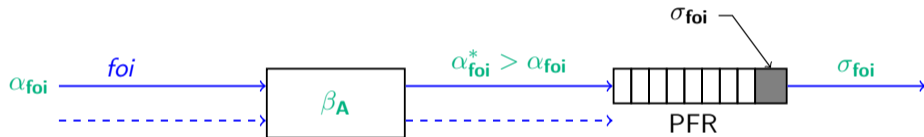
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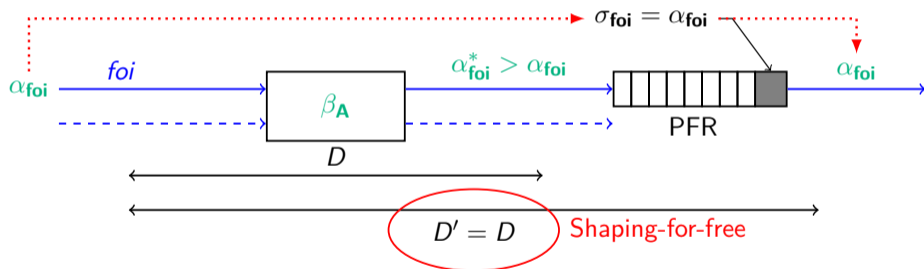
Traffic Shaping with a Per-Flow Regulator (PFR)

also called Packetized Greedy Shaper [Le Boudec and Thiran, 2001, Definition 1.7.6].



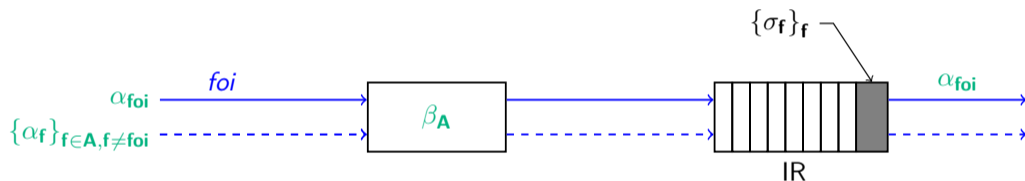
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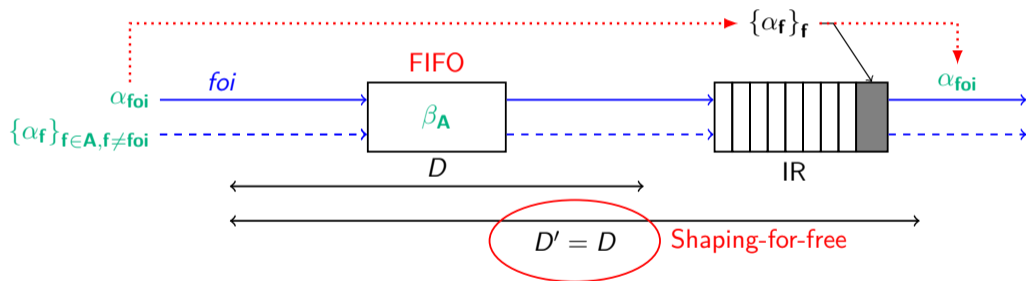
Traffic Shaping with an Interleaved Regulator (IR)

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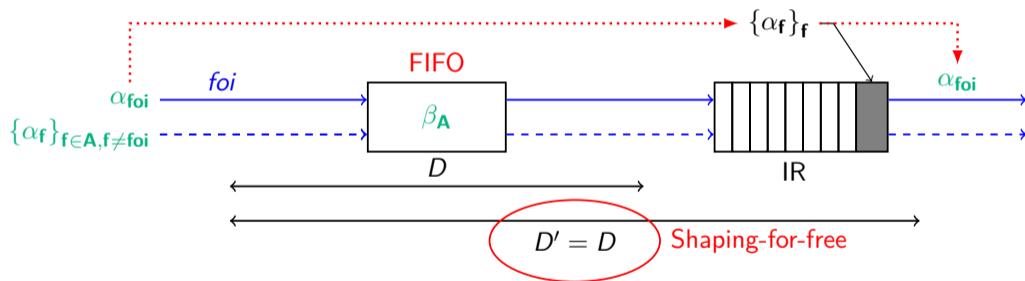
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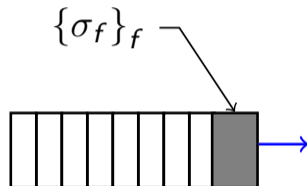
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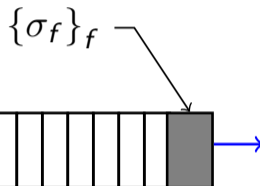


TSN Asynchronous Traffic Shaping (ATS) [IEEE, 2019] reproduces the IR in TSN.

Regulators measure **elapsed time**

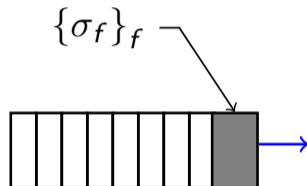


Regulators measure **elapsed time**



$$\forall f, \forall t, \forall s, \quad R^*(t + s) - R^*(t) \leq \sigma_f(s)$$

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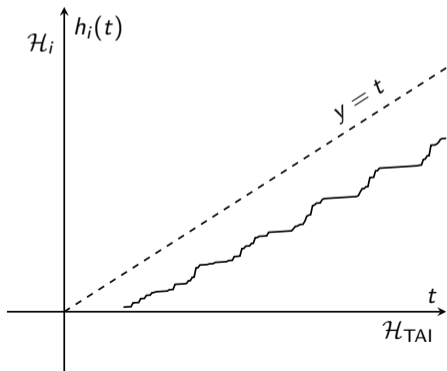
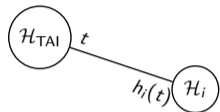
- Discussions in TSN ATS (Asynchronous Traffic Shaping) [IEEE, 2019].
- In doi:10.1145/3392145: theoretical foundations to address the problem.

Contributions

- **Time model** for $\left\{ \begin{array}{l} \text{non-synchronized} \\ \text{synchronized} \end{array} \right\}$ networks.
- A toolbox of **Network Calculus** results for $\left\{ \begin{array}{l} \text{non-synchronized} \\ \text{synchronized} \end{array} \right\}$ networks.
- Analysis of regulators $\left\{ \begin{array}{l} \text{PFR} \\ \text{IR} \end{array} \right\}$ in $\left\{ \begin{array}{l} \text{non-synchronized} \\ \text{synchronized} \end{array} \right\}$ networks.

Model for non-synchronized clocks

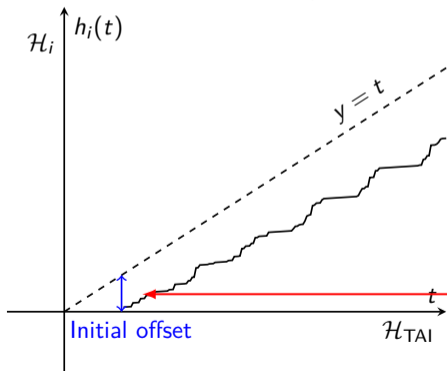
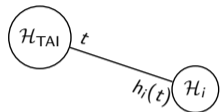
Definitions and terminology for synchronization networks [ITU, 1996]



$$h_i(t) - t = x_{i,0} + ty_{i,0} + w(t) + \psi(t)$$

Model for non-synchronized clocks

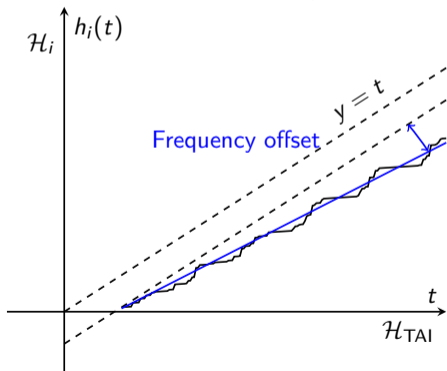
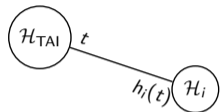
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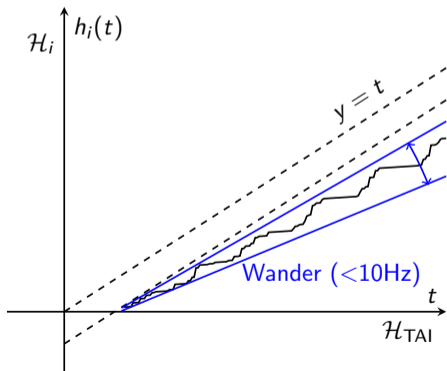
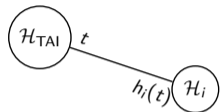


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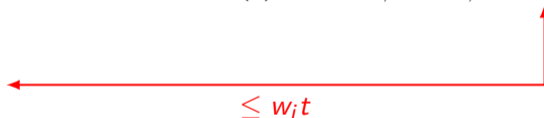
$$\leq y_{\max} t$$

Model for non-synchronized clocks

Definitions and terminology for synchronization networks [ITU, 1996]

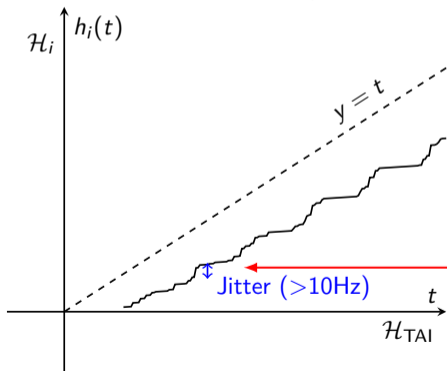
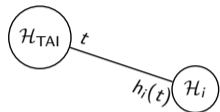


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Model for non-synchronized clocks

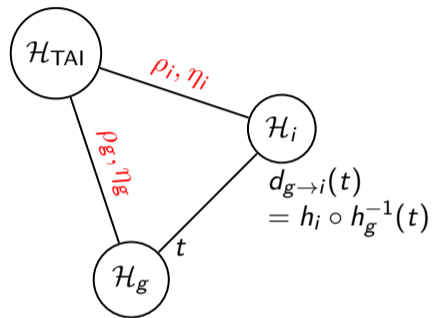
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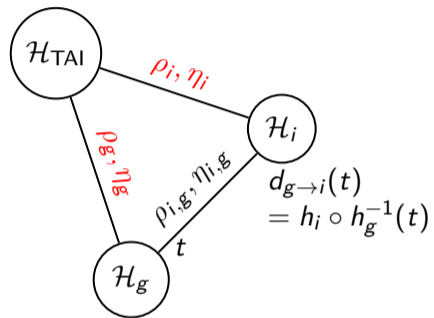
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$$\leq \eta_i$$

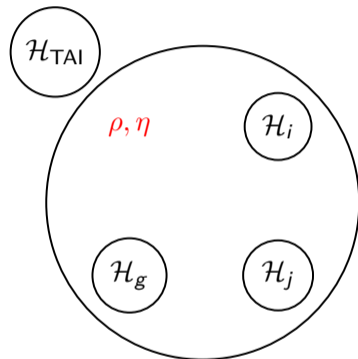
Model for non-synchronized clocks



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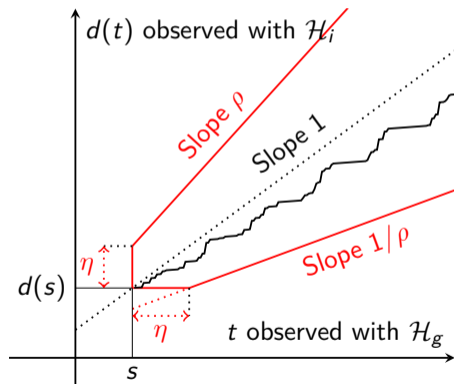
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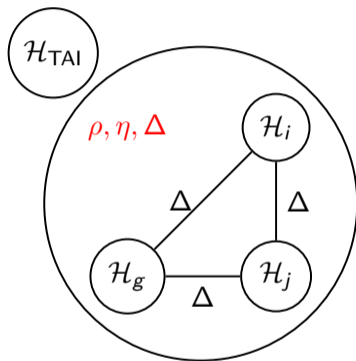
TSN: $\rho = 1 + 200ppm$
 $\eta = 4ns$

$$\forall i, g$$

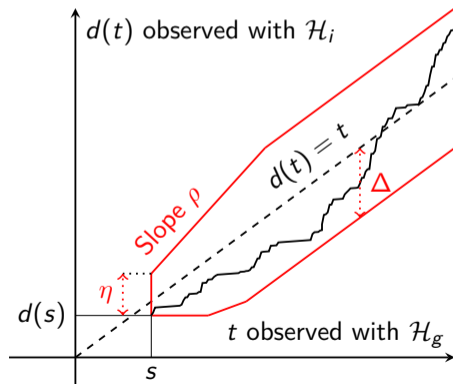
$$\frac{1}{\rho}(t - s - \eta) \leq d_{g \rightarrow i}(t) - d_{g \rightarrow i}(s) \leq (t - s)\rho + \eta$$



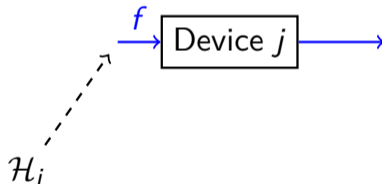
Model for synchronized clocks

TSN: $\Delta = 1\mu s$

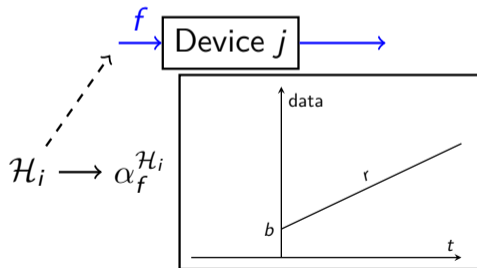
$$\forall i, g \quad \begin{cases} \frac{1}{\rho}(t - s - \eta) \leq d_{g \rightarrow i}(t) - d_{g \rightarrow i}(s) \leq (t - s)\rho + \eta \\ |d_{g \rightarrow i}(t) - t| \leq \Delta \end{cases}$$



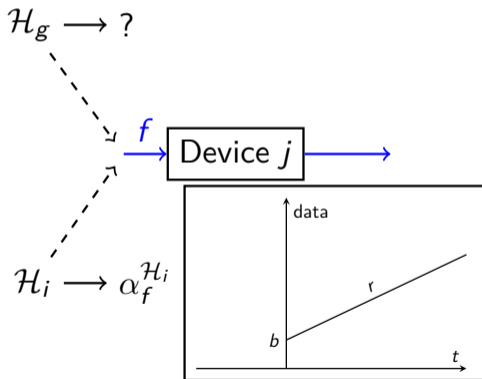
Toolbox for changing the observing clock



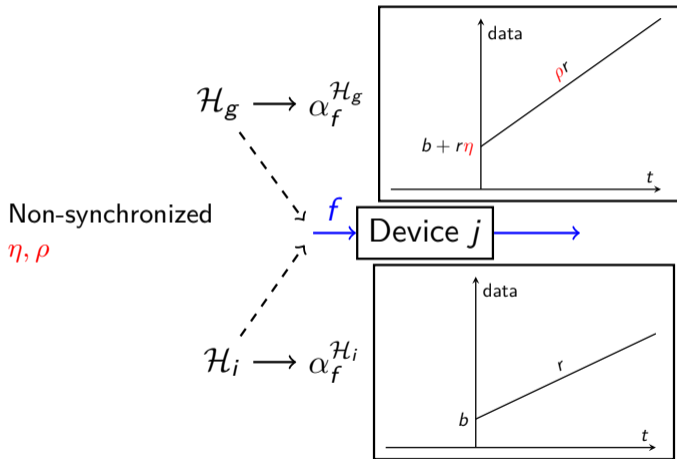
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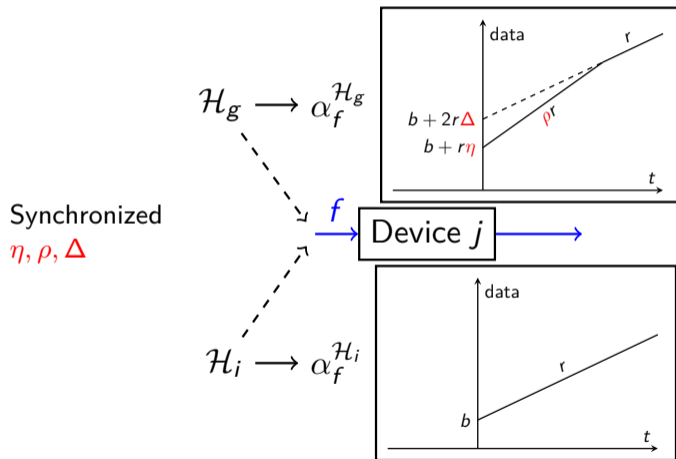
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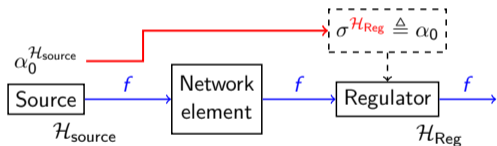


Toolbox for changing the observing clock



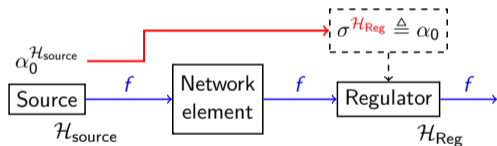
Instabilities with non-adapted regulators

Usual configuration of regulators
= Non-adapted regulator



Instabilities with non-adapted regulators

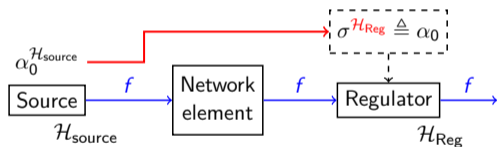
Usual configuration of regulators
= Non-adapted regulator



- Non-synchronized networks:
 - Per-flow regulator
 - Interleaved regulator
 } unstable

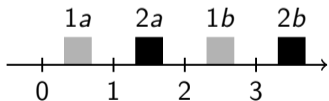
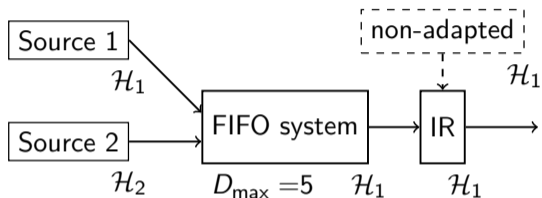
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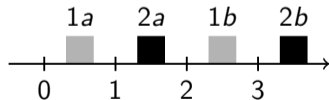
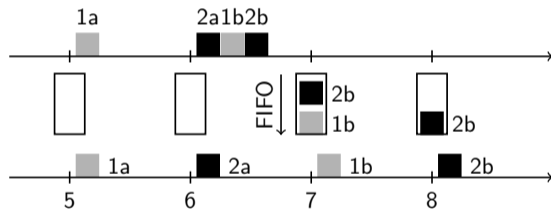
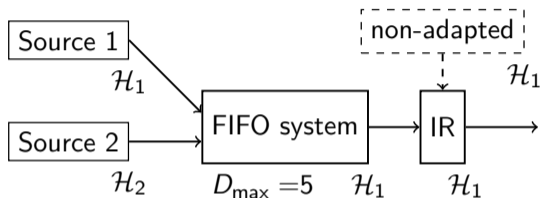


- Non-synchronized networks:
 - Per-flow regulator
 - Interleaved regulator
 } unstable
- Synchronized networks:
 - Per-flow regulator → penalty $[\Delta, 4\Delta]$
 - Interleaved regulator → **unstable** $\forall \Delta > 0$

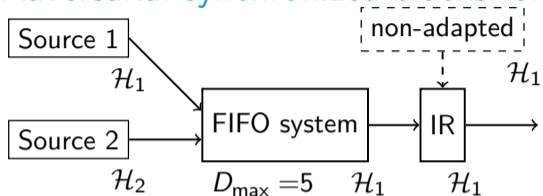
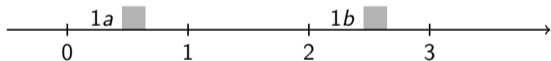
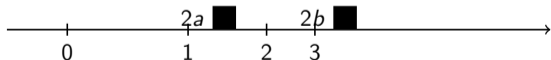
Adversarial synchronized clocks for a non-adapted IR (=ATS) (1/3)



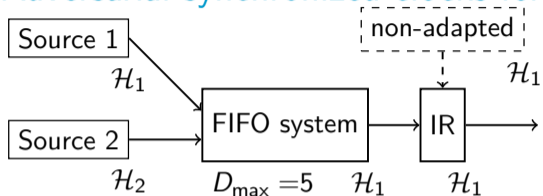
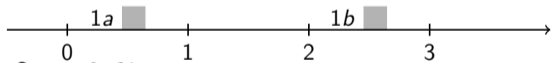
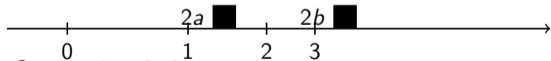
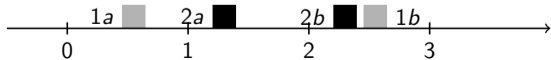
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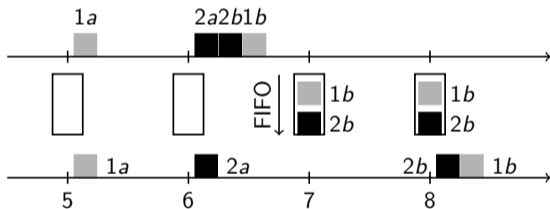
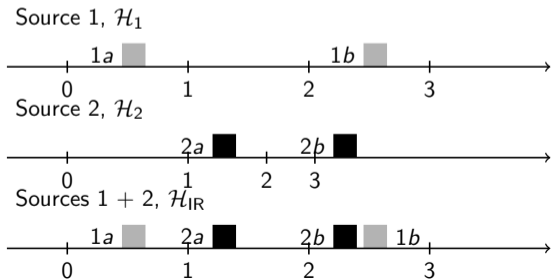
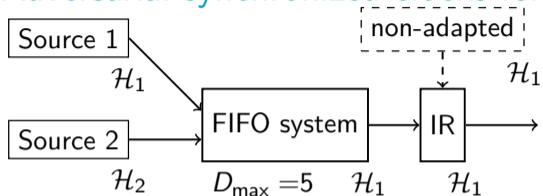
Adversarial synchronized clocks for a non-adapted IR (=ATS) (2/3)

Source 1, \mathcal{H}_1 Source 2, \mathcal{H}_2 

Adversarial synchronized clocks for a non-adapted IR (=ATS) (2/3)

Source 1, \mathcal{H}_1 Source 2, \mathcal{H}_2 Sources 1 + 2, \mathcal{H}_{IR} 

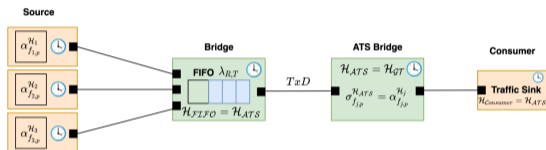
Adversarial synchronized clocks for a non-adapted IR (=ATS) (2/3)



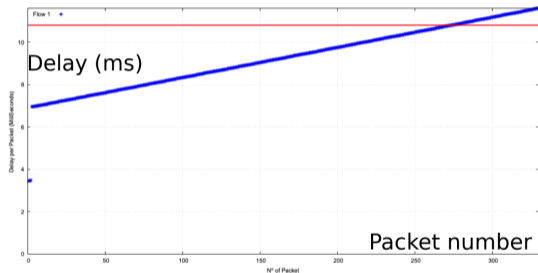
Adversarial synchronized clocks for a non-adapted IR (=ATS) (3/3)

Validation and extension through ns-3 simulations.

Example at low data rates:



- 3 sources @ 147 kB/s
- 1 queuing element @ 437.5 kB/s
- $\Delta = 1\mu s$, $\rho = 1 + 100\text{ppm}$
- using adversarial clocks
- \Rightarrow red line is Network Calculus delay bound assuming perfect clocks

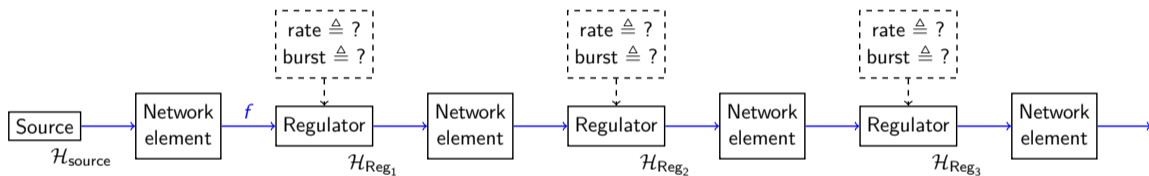


Work by Guillermo Aguirre

How to configure the regulators ?

assuming leaky-bucket-constrained flows

\forall flow f

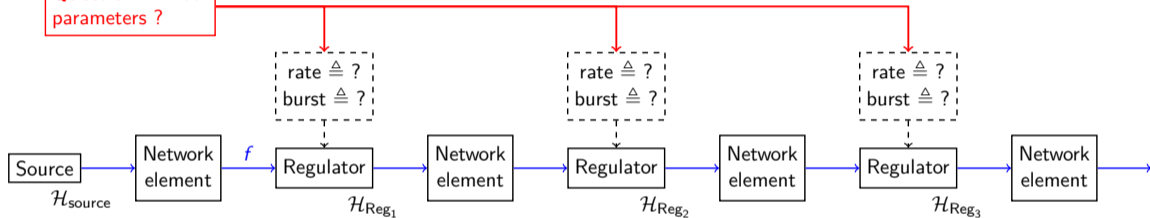


How to configure the regulators ?

assuming leaky-bucket-constrained flows

\forall flow f

Question: What parameters ?

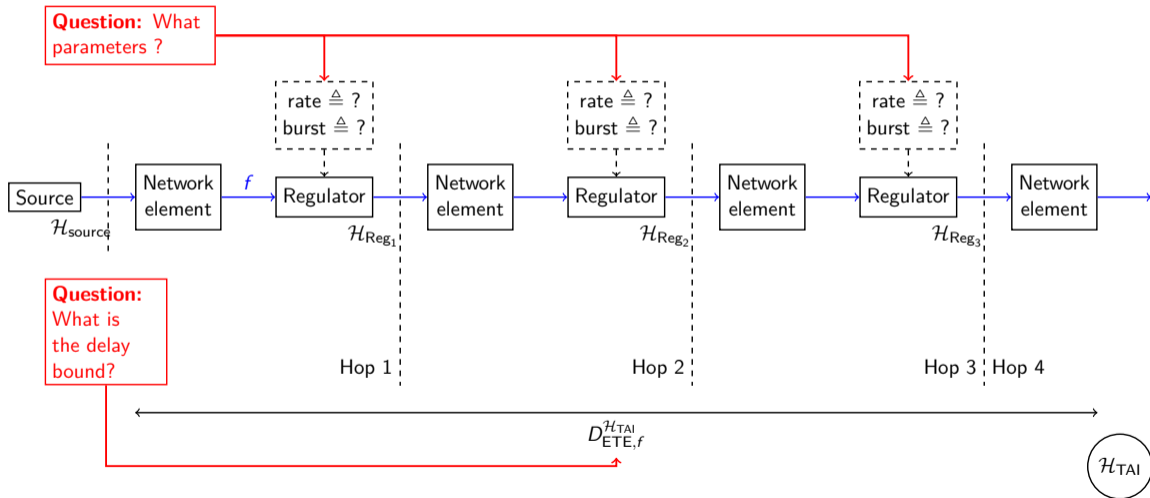


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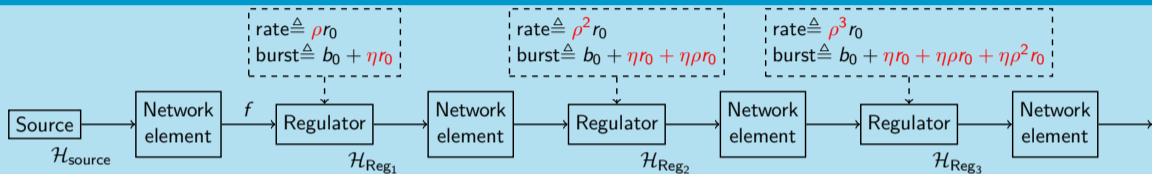
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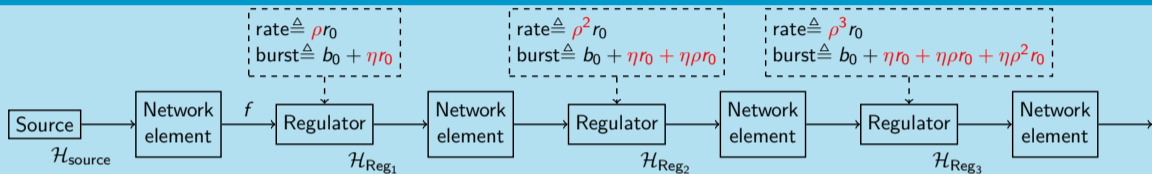
Two methods for synchronized and non-synchronized networks

Rate-and-burst cascade Works with PFR or IR

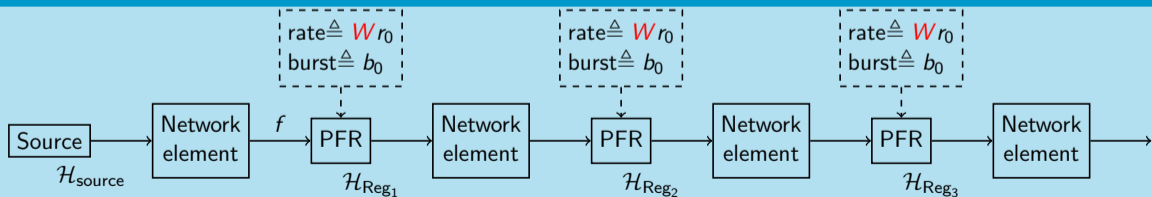


Two methods for synchronized and non-synchronized networks

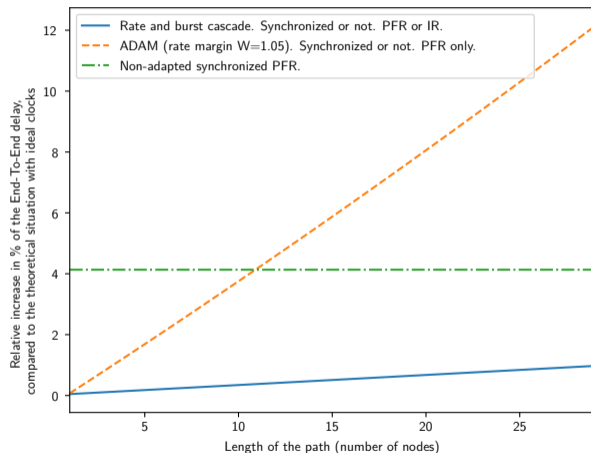
Rate-and-burst cascade Works with PFR or IR



ADAM Works with PFR only



Performance comparison



Increase of the ETE delay bound wrt ideal clocks.

Conclusion

- **Time-model** for bounding the behavior of the clocks in the network.
- **Instability** of the non-adapted ATS regulator **for any** $\Delta > 0$.
- Two methods for **configuring the regulators** in a network, relying on a **Network Calculus toolbox**.

Conclusion

- **Time-model** for bounding the behavior of the clocks in the network.
- **Instability** of the non-adapted ATS regulator **for any** $\Delta > 0$.
- Two methods for **configuring the regulators** in a network, relying on a **Network Calculus toolbox**.

Future work:

- Improvements on the ADAM method.
- Simulation of different (more realistic) clock models in ns-3.
- The toolbox could be of interest when studying other technologies / TSN components.

Bibliography I



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<https://doi.org/10.1145/1190095.1190105>.

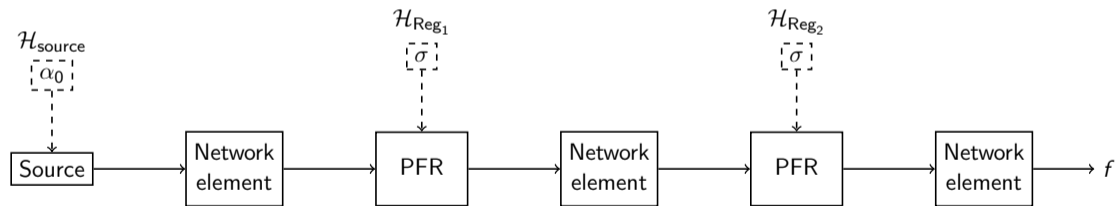


Specht, J. and Samii, S. (2016).

Urgency-Based Scheduler for Time-Sensitive Switched Ethernet Networks.

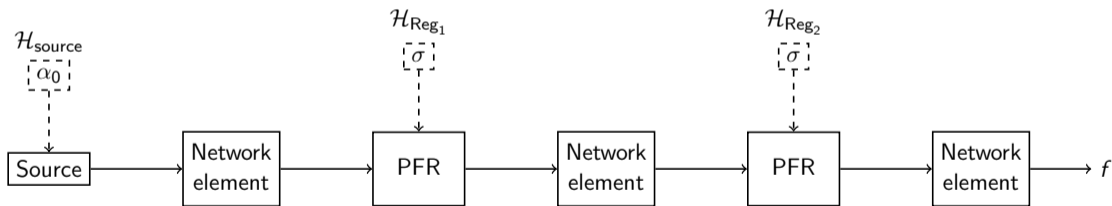
In *2016 28th Euromicro Conference on Real-Time Systems (ECRTS)*, pages 75–85.

A quick look at ADAM



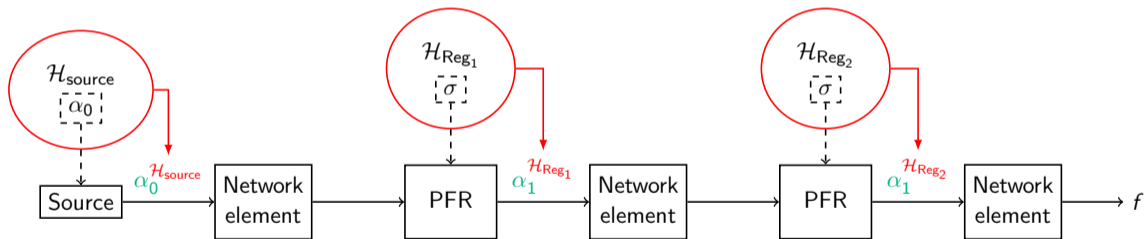
A quick look at ADAM

- Step 1 (per flow): arrival curve at each network element



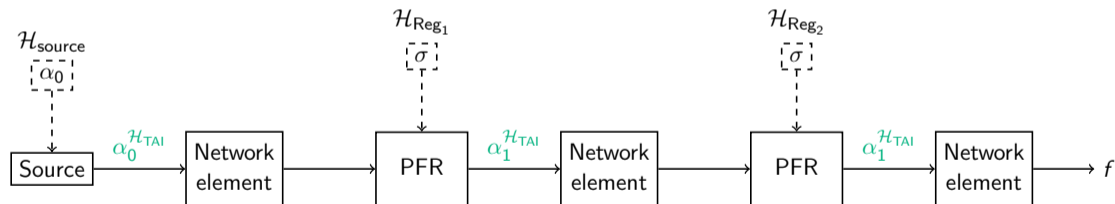
A quick look at ADAM

- Step 1 (per flow): arrival curve at each network element



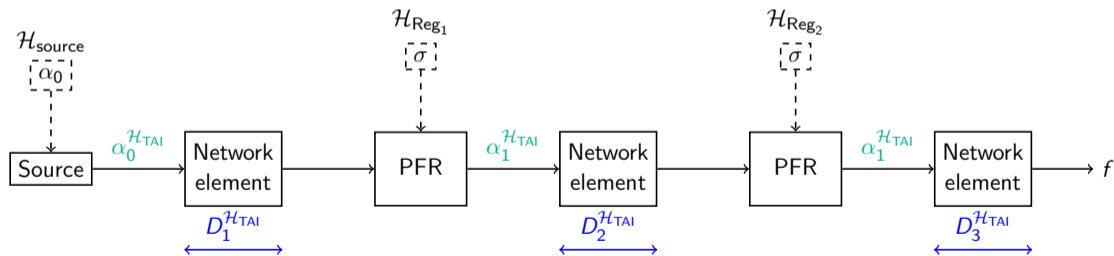
A quick look at ADAM

- Step 1 (per flow): arrival curve at each network element



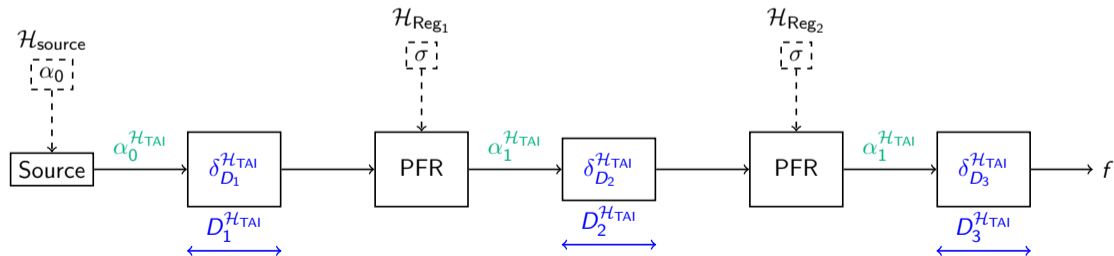
A quick look at ADAM

- Step 1 (per flow): arrival curve at each network element
- Step 2 (per node): compute delay bounds



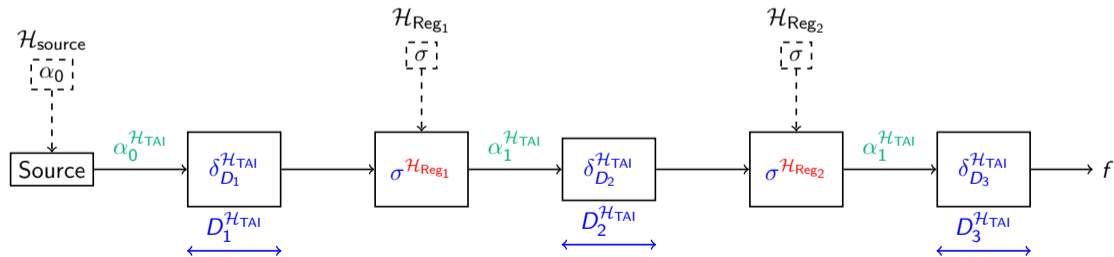
A quick look at ADAM

- Step 1 (per flow): arrival curve at each network element
- Step 2 (per node): compute delay bounds
- Step 3 (per flow): service-curve characterization and end-to-end delay bound



A quick look at ADAM

- Step 1 (per flow): arrival curve at each network element
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A quick look at ADAM

- Step 1 (per flow): arrival curve at each network element
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